

## Chapter 10

### Exercise 10.1

(a), (b), (c), (e), (f), (h) are true; (d), (g) are false.

### Exercise 10.2

Use [wolfgangschwarz.net/trees/](http://wolfgangschwarz.net/trees/). Note that the website uses slightly different identity rules: instead of the Self-Identity rule, it has a rule for closing any branch that contains a statement of the form  $\tau \neq \tau$ .

### Exercise 10.3

- (a)  $W = \{w\}$ ,  $wRw$ ,  $D = \{\text{Alice}\}$ ,  $V(F, w) = \{\text{Alice}\}$ ,  $V(G, w) = \emptyset$
- (b)  $W = \{w, v\}$ ,  $wRw$  and  $wRv$ ,  $D = \{\text{Alice}, \text{Bob}\}$ ,  $V(F, w) = \{\text{Alice}\}$ ,  $V(F, v) = \{\text{Bob}\}$
- (c)  $W = \{w, v\}$ ,  $wRw$  and  $wRv$ ,  $D = \{\text{Alice}, \text{Bob}\}$ ,  $V(F, w) = \{\text{Alice}\}$ ,  $V(F, v) = \emptyset$
- (d)  $W = \{w, v\}$ ,  $wRw$  and  $wRv$ ,  $D = \{\text{Alice}, \text{Bob}\}$ ,  $V(P, w) = \{\text{Alice}\}$ ,  $V(P, v) = \emptyset$ ,  $V(Q, w) = \{\text{Alice}\}$ ,  $V(Q, v) = \emptyset$

### Exercise 10.4

$\Box \forall x \exists y (x = y) \rightarrow \forall x \Box \exists y (x = y)$  is an instance of the Converse Barcan Formula. If we read the box as a relevant kind of circumstantial necessity, and Loafy could have failed to exist, the consequent of this conditional is false. But the antecedent is true.

### Exercise 10.5

(1) is equivalent to the Barcan Formula, (4) to the Converse Barcan Formula. (2) is highly implausible. (1) and (4) are often regarded as implausible, for the reasons I discuss in the text. (3) is about as plausible or implausible as the Converse Barcan Formula.

### Exercise 10.6

- (a)
- |  |    |   |     |        |
|--|----|---|-----|--------|
|  | 1. | $\exists x \Box Fx \rightarrow \Box \exists x Fx$ | (w) | (Ass.) |
|  | 2. | $\exists x \Box Fx$                               | (w) | (1)    |
|  | 3. | $\neg \Box \exists x Fx$                          | (w) | (1)    |
|  | 4. | $\Box Fa$   | (w) | (2)    |
|  | 5. | $wRv$   |     | (3)    |
|  | 6. | $\neg \exists x Fx$                               | (v) | (3)    |
|  | 7. | $Fa$  | (v) | (4,5)  |
|  | 8. | $a = a$   | (v) | (7)    |
- 
- |    |            |         |  |                      |
|----|------------|---------|--|----------------------|
|    |            |         |  |                      |
| 9. | $a \neq a$ | (v) (6) |  | 9. $\neg Fa$ (v) (6) |
|    | x          |         |  | x                    |

(b) DIY. The tree has four branches. I can't typeset it.

- (c)
- |  |    |                             |     |        |
|--|----|-----------------------------|-----|--------|
|  | 1. | $\neg \Box \exists x x = x$ | (w) | (Ass.) |
|  | 2. | $wRv$                       |     | (1)    |
|  | 3. | $\neg \exists x x = x$      | (v) | (1)    |
|  | 4. | $a = a$                     | (v) | (Ex.)  |
- 
- |    |            |         |  |                       |
|----|------------|---------|--|-----------------------|
|    |            |         |  |                       |
| 9. | $a \neq a$ | (v) (3) |  | 9. $a \neq a$ (v) (3) |
|    | x          |         |  | x                     |

- (d)
- |  |    |   |     |        |
|--|----|---|-----|--------|
|  | 1. | $\neg(\Diamond Fa \rightarrow \Diamond \exists x Fx)$ | (w) | (Ass.) |
|  | 2. | $\Diamond Fa$   | (w) | (1)    |
|  | 3. | $\neg \Diamond \exists x Fx$                          | (w) | (1)    |
|  | 4. | $wRv$   |     | (2)    |
|  | 5. | $Fa$  | (v) | (2)    |
|  | 6. | $a = a$   | (v) | (5)    |
|  | 7. | $\neg \exists x Fx$                                   | (v) | (3,4)  |
- 
- |    |            |         |  |                       |
|----|------------|---------|--|-----------------------|
|    |            |         |  |                       |
| 9. | $a \neq a$ | (v) (3) |  | 10. $\neg Fa$ (v) (3) |
|    | x          |         |  | x                     |

|        |   |            |
|--------|---|------------|
| (e) 1. | $\neg(a=b \rightarrow \Box(a=a \rightarrow a=b))$ | (w) (Ass.) |
| 2.     | $a=b$   | (w) (1)    |
| 3.     | $\neg\Box(a=a \rightarrow a=b)$                   | (w) (1)    |
| 4.     | $wRv$   | (3)        |
| 5.     | $\neg(a=a \rightarrow a=b)$                       | (v) (3)    |
| 6.     | $a=a$   | (v) (5)    |
| 7.     | $\neg a=b$  | (v) (5)    |
| 8.     | $a=b$   | (v) (2,6)  |
|        | x   |            |

**Exercise 10.7**

In the definition of a model, we could allow the interpretation function to be undefined for some names. We might also allow the sets  $D_w$  to be empty. We could leave the truth definition as it is.

**Exercise 10.8**

In the Superman case, Clark Kent and Superman are the same person, but Lois Lane doesn't know that they are. So we appear to have  $s=c$  but not  $\Box(s=c)$ . Similarly, in the Julius case, Julius and Whitcomb L. Judson are the same person, but one may well not know that they are. In the Goliath case, we have  $\text{Lumpl} = \text{Goliath}$  without it being metaphysically necessary that  $\text{Lumpl} = \text{Goliath}$ , as there are worlds in which  $\text{Lumpl}$  is a bowl and  $\text{Goliath}$  is not.

**Exercise 10.9**

We would assume that (i) the name  $g$  picks out a statue at all accessible worlds, (ii)  $l$  picks out a lump of clay at all accessible worlds, and (iii) at the actual world,  $l$  and  $g$  pick out the same thing: the statue-shaped lump on the shelf.

**Exercise 10.10**

The premises are  $\Box\exists x(x = i)$  and  $\neg\Box\exists x(x = b)$ . The conclusion is  $i \neq b$ . The

argument is CK-valid and VK-valid.

**Exercise 10.11**

Translation:  $\exists x(Tx \wedge Wx \wedge \neg K Wx \wedge \neg K \neg Wx)$ , where  $T$  translates ‘– is a ticket’ and ‘– will win’.

If variables are directly referential, then this sentence is true in any scenario in which I don’t know which ticket will win.

**Exercise 10.12**

To render  $\forall x \forall y (x = y \rightarrow \Box x = y)$  valid, we can restrict the eligible individual concepts in a model as follows. For any individual concepts  $f$  and  $g$  and worlds  $w$  and  $v$ , if  $w R v$  and  $f(w) = g(w)$  then  $f(v) = g(w)$ . (We do not stipulate that if  $w R v$  and  $f(v) = g(v)$  then  $f(w) = g(w)$ , which would render the necessity of distinctness valid.)