Logic 2: Modal Logic

Lecture 6

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The system T

The system T

Schema		S5	K
K	$\Box(A \to B) \to (\Box A \to \Box B)$	✓	✓
T	$\Box A \to A$	✓	—
D	$\Box A \to \Diamond A$	✓	—
B	$A \to \Box\Diamond A$	✓	—
4	$\Box A \to \Box\Box A$	✓	—
5	$\Diamond A \to \Box\Diamond A$	✓	—
G	$\Diamond\Box A \to \Box\Diamond A$	✓	—

The system T

Schema		S5	K	T
K	$\Box(A \to B) \to (\Box A \to \Box B)$	✓	✓	✓
T	$\Box A \to A$	✓	—	✓
D	$\Box A \to \Diamond A$	✓	—	✓
B	$A \to \Box\Diamond A$	✓	—	—
4	$\Box A \to \Box\Box A$	✓	—	—
5	$\Diamond A \to \Box\Diamond A$	✓	—	—
G	$\Diamond\Box A \to \Box\Diamond A$	✓	—	—

The system T

Let $\Box p$ mean ‘you know that p’.

- $\Box p \rightarrow p$ is plausibly valid.
- $\Box p \rightarrow \Diamond p$ is plausibly valid.
- $p \rightarrow \Box\Diamond p$ is not.

Suppose you falsely believe $\neg p$.

- p is true.
- You believe that you know $\neg p$.
- You don’t believe that you don’t know $\neg p$.
- So you don’t know that you don’t know $\neg p$.
- $\neg\Box\neg\Box\neg p$
- $\neg\Box\Diamond p$

The system T

Kripke model

A Kripke model of $\mathfrak{L}_M$ is a triple $\langle W, R, V \rangle$ consisting of

- a non-empty set W,
- a binary relation R on W, and
- a function V that assigns to each sentence letter of $\mathfrak{L}_M$ and each element of W a truth-value.

Intuitively: wRv iff v is possible relative to w.

Intuitively: wRv iff v is compatible with your knowledge at w.

The system T

Whatever you know at a world is true at that world.

So every world w is compatible with your knowledge at w.

So every world is accessible from itself.

The system T

Whatever you know at a world is true at that world.

So every world *w* is compatible with your knowledge at *w*.

So every world is accessible from itself.

The system T

Definition: K-valid

A sentence A is **K-valid** (for short, $\models_K A$) iff A is true at every world in every Kripke model.

Definition: T-valid

A sentence A is **T-valid** (for short, $\models_T A$) iff A is true at every world in every Kripke model in which each world has access to itself.

The set of all T-valid sentences is **system T**.

The system T

Schema		S5	K	T
K	$\Box(A \to B) \to (\Box A \to \Box B)$	✓	✓	✓
T	$\Box A \to A$	✓	—	✓
D	$\Box A \to \Diamond A$	✓	—	✓
B	$A \to \Box\Diamond A$	✓	—	—
4	$\Box A \to \Box\Box A$	✓	—	—
5	$\Diamond A \to \Box\Diamond A$	✓	—	—
G	$\Diamond\Box A \to \Box\Diamond A$	✓	—	—

The system T

New tree rule for system T:

$$\vdots$$

$$\omega R \omega$$

↑

old

The system T

Target sentence: $\Box p \rightarrow p$

1.	$\neg(\Box p \rightarrow p)$	(w)	(Ass.)
2.	$\Box p$	(w)	(1)
3.	$\neg p$	(w)	(1)
4.	wRw		(Ref.)
5.	p	(w)	(2,4)
	x		

The system S4

The system S4

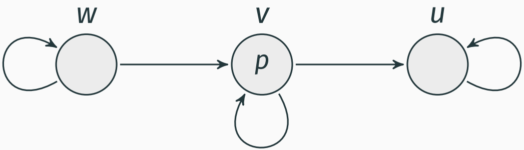
Let $\Box p$ mean 'you know that p'.

- $\Box p \rightarrow p$ is plausibly valid.
- $\Box p \rightarrow \Diamond p$ is plausibly valid.
- $p \rightarrow \Box\Diamond p$ is not.
- $\Box p \rightarrow \Box\Box p$ might be valid.

The system S4

Schema		S5	K	T	S4
K	$\Box(A \rightarrow B) \rightarrow (\Box A \rightarrow \Box B)$	✓	✓	✓	✓
T	$\Box A \rightarrow A$	✓	—	✓	✓
D	$\Box A \rightarrow \Diamond A$	✓	—	✓	✓
B	$A \rightarrow \Box\Diamond A$	✓	—	—	—
4	$\Box A \rightarrow \Box\Box A$	✓	—	—	✓
5	$\Diamond A \rightarrow \Box\Diamond A$	✓	—	—	—
G	$\Diamond\Box A \rightarrow \Box\Diamond A$	✓	—	—	—

The system S4

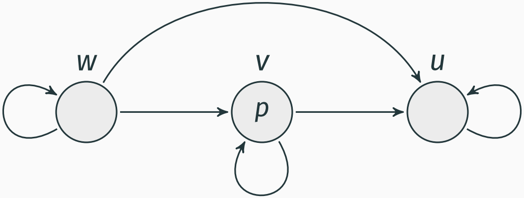


$\Box A \to \Box\Box A$ is valid iff $\Diamond\Diamond A \to \Diamond A$ is valid.

$\Diamond\Diamond A \to \Diamond A$:

- "If you can get somewhere in two steps, you can get there in one step."
- If v is compatible with your knowledge at w, and u is compatible with your knowledge at v, then u is compatible with your knowledge at w.

The system S4



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- "If you can get somewhere in two steps, you can get there in one step."
- If v is compatible with your knowledge at w, and u is compatible with your knowledge at v, then u is compatible with your knowledge at w.

The system S4

Definition: S4-valid

A sentence A is **S4-valid** (for short, $\models_{S4} A$) iff A is true at every world in every Kripke model in which

(a) each world has access to itself, and

(b) if a world w has access to some world v, and v has access to a world u, then w has access to u.

The set of all S4-valid sentences is **system S4**.

The system S4

Schema		S5	K	T	S4
K	$\Box(A \to B) \to (\Box A \to \Box B)$	✓	✓	✓	✓
T	$\Box A \to A$	✓	—	✓	✓
D	$\Box A \to \Diamond A$	✓	—	✓	✓
B	$A \to \Box\Diamond A$	✓	—	—	—
4	$\Box A \to \Box\Box A$	✓	—	—	✓
5	$\Diamond A \to \Box\Diamond A$	✓	—	—	—
G	$\Diamond\Box A \to \Box\Diamond A$	✓	—	—	—

The system S4

New tree rule for system S4:

$$\begin{array}{c} \omega R \nu \\ \nu R \upsilon \\ \vdots \\ \omega R \upsilon \end{array}$$

The system S4

Target sentence: $\Box A \to \Box\Box A$

1.	$\neg(\Box A \to \Box\Box A)$	(w)	(Ass.)
2.	$\Box A$	(w)	(1)
3.	$\neg\Box\Box A$	(w)	(1)
4.	wRv		(3)
5.	$\neg\Box A$	(v)	(3)
6.	vRu		(5)
7.	$\neg A$	(u)	(5)
8.	wRu		(4,6,Tr.)
5.	A	(u)	(2,8)
	X		

Further systems

Further systems

A relation R on a set W is called **reflexive** if every member of W is R-related to itself.

Definition: T-valid

A sentence A is **T-valid** (for short, $\models_T A$) iff A is true at every world in every Kripke model in which each world has access to itself.

Definition: T-valid

A sentence A is **T-valid** (for short, $\models_T A$) iff A is true at every world in every Kripke model with a reflexive accessibility relation.

A relation R on a set W is called **transitive** if whenever wRv and vRu, then wRu.

Definition: S4-valid

A sentence A is **S4-valid** (for short, $\vDash_{S4} A$) iff A is true at every world in every Kripke model in which

(a) each world has access to itself, and

(b) if a world w has access to some world v, and v has access to a world u, then w has access to u.

Definition: S4-valid

A sentence A is **S4-valid** (for short, $\vDash_{S4} A$) iff A is true at every world in every Kripke model whose accessibility relation is reflexive and transitive.

Further systems

Definition: K4-valid

A sentence A is **K4-valid** (for short, $\models_{K4} A$) iff A is true at every world in every Kripke model whose accessibility relation is transitive.

Further systems

Schema		S5	K	T	S4	K4
K	$\Box(A \to B) \to (\Box A \to \Box B)$	✓	✓	✓	✓	✓
T	$\Box A \to A$	✓	—	✓	✓	—
D	$\Box A \to \Diamond A$	✓	—	✓	✓	—
B	$A \to \Box\Diamond A$	✓	—	—	—	—
4	$\Box A \to \Box\Box A$	✓	—	—	✓	✓
5	$\Diamond A \to \Box\Diamond A$	✓	—	—	—	—
G	$\Diamond\Box A \to \Box\Diamond A$	✓	—	—	—	—

Further systems

A relation R on a set W is called **serial** if every member of W is R-related to some member of W. ("No dead ends")

Further systems

A relation R on a set W is called **serial** if every member of W is R-related to some member of W. ("No dead ends")

Definition: D-valid

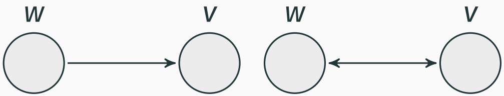
A sentence A is **D-valid** (for short, $\models_D A$) iff A is true at every world in every Kripke model whose accessibility relation is serial.

Further systems

Schema		S5	K	T	S4	K4	D
K	$\Box(A \to B) \to (\Box A \to \Box B)$	✓	✓	✓	✓	✓	✓
T	$\Box A \to A$	✓	—	✓	✓	—	—
D	$\Box A \to \Diamond A$	✓	—	✓	✓	—	✓
B	$A \to \Box\Diamond A$	✓	—	—	—	—	—
4	$\Box A \to \Box\Box A$	✓	—	—	✓	✓	—
5	$\Diamond A \to \Box\Diamond A$	✓	—	—	—	—	—
G	$\Diamond\Box A \to \Box\Diamond A$	✓	—	—	—	—	—

Further systems

A relation R on a set W is called **symmetric** if whenever wRv then vRw.



Further systems

A relation R on a set W is called **symmetric** if whenever wRv then vRw.

Definition: B-valid

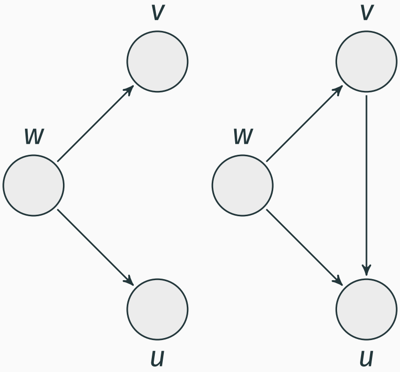
A sentence A is **B-valid** (for short, $\models_B A$) iff A is true at every world in every Kripke model whose accessibility relation is reflexive and symmetric.

Further systems

Schema		S5	K	T	S4	K4	D	B
K	$\Box(A \to B) \to (\Box A \to \Box B)$	✓	✓	✓	✓	✓	✓	✓
T	$\Box A \to A$	✓	—	✓	✓	—	—	✓
D	$\Box A \to \Diamond A$	✓	—	✓	✓	—	✓	✓
B	$A \to \Box\Diamond A$	✓	—	—	—	—	—	✓
4	$\Box A \to \Box\Box A$	✓	—	—	✓	✓	—	—
5	$\Diamond A \to \Box\Diamond A$	✓	—	—	—	—	—	—
G	$\Diamond\Box A \to \Box\Diamond A$	✓	—	—	—	—	—	—

Further systems

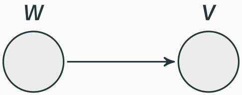
A relation R on a set W is called **euclidean** if whenever wRv and wRu, then vRu.

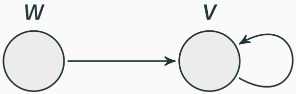


Further systems

A relation R is called **euclidean** if whenever wRv and wRu, then vRu.

Note: w, v, and u need not be distinct worlds! $\forall x \forall y \forall z(xRy \wedge xRz \rightarrow yRz)$.





Further systems

A relation R on a set W is called **euclidean** if whenever wRv and wRu, then vRu.

Definition

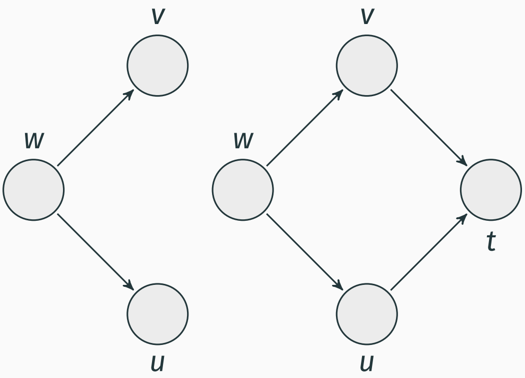
A sentence A is **K45-valid** (for short, $\models_{K45} A$) iff A is true at every world in every Kripke model whose accessibility relation is transitive and euclidean.

Further systems

Schema		S5	K	T	S4	K4	D	B	K45
K	$\Box(A \rightarrow B) \rightarrow (\Box A \rightarrow \Box B)$	✓	✓	✓	✓	✓	✓	✓	✓
T	$\Box A \rightarrow A$	✓	—	✓	✓	—	—	✓	—
D	$\Box A \rightarrow \Diamond A$	✓	—	✓	✓	—	✓	✓	—
B	$A \rightarrow \Box\Diamond A$	✓	—	—	—	—	—	✓	—
4	$\Box A \rightarrow \Box\Box A$	✓	—	—	✓	✓	—	—	✓
5	$\Diamond A \rightarrow \Box\Diamond A$	✓	—	—	—	—	—	—	✓
G	$\Diamond\Box A \rightarrow \Box\Diamond A$	✓	—	—	—	—	—	—	—

Further systems

A relation R on a set W is called **convergent** if whenever wRv and wRu, then there is some t such that wRt and vRt.



A relation R on a set W is called **convergent** if whenever wRv and wRu, then there is some t such that wRt and vRt.

Definition

A sentence A is **S4.2-valid** (for short, $\vDash_{S4.2} A$) iff A is true at every world in every Kripke model whose accessibility relation is reflexive, transitive, and convergent.

Further systems

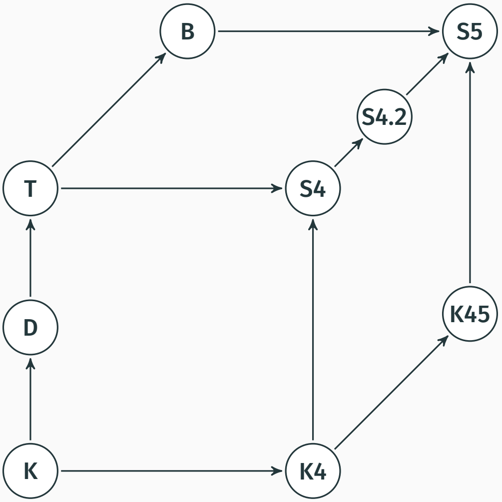
Schema		S5	K	T	S4	K4	D	B	K45	S4.2
K	$\Box(A \to B) \to (\Box A \to \Box B)$	✓	✓	✓	✓	✓	✓	✓	✓	✓
T	$\Box A \to A$	✓	—	✓	✓	—	—	✓	—	✓
D	$\Box A \to \Diamond A$	✓	—	✓	✓	—	✓	✓	—	✓
B	$A \to \Box\Diamond A$	✓	—	—	—	—	—	✓	—	—
4	$\Box A \to \Box\Box A$	✓	—	—	✓	✓	—	—	✓	✓
5	$\Diamond A \to \Box\Diamond A$	✓	—	—	—	—	—	—	✓	—
G	$\Diamond\Box A \to \Box\Diamond A$	✓	—	—	—	—	—	—	—	✓

Further systems

Schema		Condition On R
T	$\Box A \to A$	R is reflexive: every world in W is accessible from itself
D	$\Box A \to \Diamond A$	R is serial: every world in W can access some world in W
B	$A \to \Box\Diamond A$	R is symmetric: whenever wRv then vRw
4	$A \to \Box\Box A$	R is transitive: whenever wRv and vRu, then wRu
5	$\Diamond A \to \Box\Diamond A$	R is euclidean: whenever wRv and wRu, then vRu
G	$\Diamond\Box A \to \Box\Diamond A$	R is convergent: whenever wRv and wRu, then there is some t such that vRt and uRt

Further systems

K	–
T	R is reflexive
D	R is serial
K4	R is transitive
K45	R is transitive and euclidean
B	R is reflexive and symmetric
S4	R is reflexive and transitive
S4.2	R is reflexive, transitive, and convergent
S5	R is reflexive, transitive, and symmetric
S5	R is universal



Further systems