Logic 2: Modal Logic

Lecture 13

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28 October 2019

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Fatalism

Fatalism

Operators

$\mathrm{P}\,A$	it was once the case that A
$\mathrm{F}\,A$	it will once be the case that A
$\Diamond A$	you can bring it about that A

Assumptions

(i) $\mathrm{F}\,A \models \mathrm{P}\,\mathrm{F}\,A$

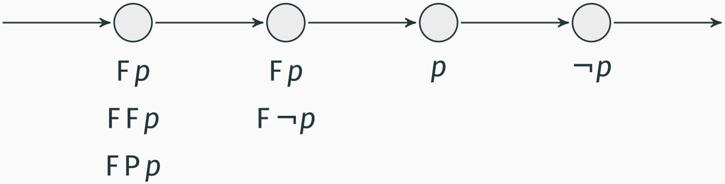
(ii) $\neg\,\mathrm{P}\,A \models \neg\Diamond\,\mathrm{P}\,A$

(iii) If $A \models B$ then $\Diamond A \models \Diamond B$

1. Suppose q neither is, nor was, nor will ever be the case.
2. So $\neg\,\mathrm{P}\,\mathrm{F}\,q$.
3. Then $\neg\Diamond\,\mathrm{P}\,\mathrm{F}\,q$, by (ii).
4. But Fq entails $\mathrm{P}\,\mathrm{F}\,q$, by (i).
5. So $\Diamond Fq$ entails $\Diamond\,\mathrm{P}\,\mathrm{F}\,q$, by (iii).
6. So $\neg\Diamond\,\mathrm{F}\,q$, by (3), (5), and Modus Tollens.

Temporal Models

Temporal Models



Once the structure of time and the truth-values of sentence letters relative to each time is fixed, we can compute the truth-value of every sentence at every time.

Temporal Models

Temporal Model

A **temporal model** consists of

- a non-empty set T (of "times"),
- a binary relation $<$ on T (the **precedence relation**),
- a function V that assigns to each sentence letter of $\mathfrak{L}_T$ and each member of T a truth-value (1 or 0).

Temporal Models

Standard Temporal Semantics

(a)	$M, t \models p$	iff $V(p, t) = 1$.
(b)	$M, t \models \neg A$	iff $M, t \not\models A$.
(c)	$M, t \models A \wedge B$	iff $M, t \models A$ and $M, t \models B$.
(d)	$M, t \models A \vee B$	iff $M, t \models A$ or $M, t \models B$.
(e)	$M, t \models A \rightarrow B$	iff $M, t \models B$ or $M, t \not\models A$.
(f)	$M, t \models A \leftrightarrow B$	iff $M, t \models (A \rightarrow B)$ and $M, t \models (B \rightarrow A)$.
(g)	$M, t \models \mathrm{F}A$	iff $M, s \models A$ for some $s \in T$ such that $t < s$.
(h)	$M, t \models \mathrm{G}A$	iff $M, s \models A$ for all $s \in T$ such that $t < s$.
(i)	$M, t \models \mathrm{P}A$	iff $M, s \models A$ for some $s \in T$ such that $s < t$.
(j)	$M, t \models \mathrm{H}A$	iff $M, s \models A$ for all $s \in T$ such that $s < t$.

Temporal Models

Some consequences:

- $\models \mathsf{G}(A \to B) \to (\mathsf{G}\,A \to \mathsf{G}\,B)$.
- $\models \mathsf{G}(A \land B) \leftrightarrow (\mathsf{G}\,A \land \mathsf{G}\,B)$.
- If $\models A$ then $\models \mathsf{G}\,A$.
- If $A \models B$ then $\mathsf{G}\,A \models \mathsf{G}\,B$.

A potential problem

I have always lived either in Brimbank or Hobsons Bay.

$\rightsquigarrow$ I have lived in Brimbank.

$\rightsquigarrow$ I have lived in Hobsons Bay.

If $\mathrm{H}(p \vee q)$ entails $\mathrm{P}\,q$, then $\mathrm{H}(p)$ can't entail $\mathrm{H}(p \vee q)$.

The flow of time

The flow of time

$t < r$ iff some path along the arrows leads from t to r.

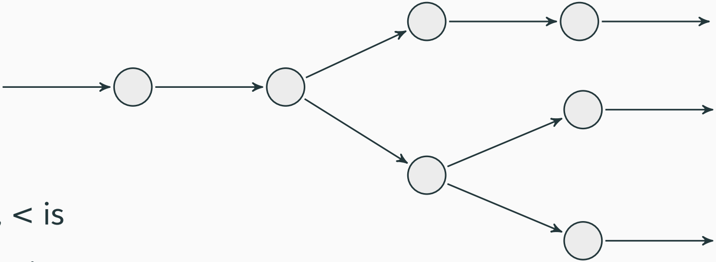
In this frame, $<$ is

- transitive: if $t < s$ and $s < r$ then $t < r$.
- asymmetric: if $t < s$ then $s \nless t$.
- irreflexive: for no time, $t < t$.
- discrete: if $t < s$ then there is an r such that $t < r$ and for no x, $t < x < r$.
- connected: either $t < s$ or $t = s$ or $s < t$.

The flow of time

- Transitivity corresponds to **4**: $\mathsf{G}A \rightarrow \mathsf{G}\mathsf{G}A$.
- Asymmetry corresponds to nothing.
- Irreflexivity corresponds to nothing.
- Discreteness corresponds to $(\mathsf{F}(p \vee \neg p) \wedge A \wedge \mathsf{H}A) \rightarrow \mathsf{F}\mathsf{G}A$.
- Connectedness corresponds to
  $(\mathsf{F}\mathsf{P}A \rightarrow (\mathsf{F}A \vee A \vee \mathsf{P}A)) \wedge (\mathsf{P}\mathsf{F}A \rightarrow (\mathsf{P}A \vee A \vee \mathsf{F}A))$

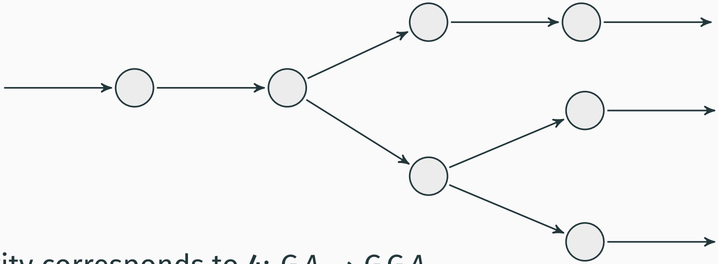
The flow of time



In this frame, $<$ is

- transitive: if $t < s$ and $s < r$ then $t < r$.
- asymmetric: if $t < s$ then $s \not< t$.
- discrete: if $t < s$ then there is an r such that $t < r$ and for no x, $t < x < r$.
- left-linear: if $t < r$ and $s < r$ then either $t < s$ or $t = s$ or $s < t$.
- right-branching: for some $r < t$ and $r < s$, neither $t < s$ or $t = s$ or $s < t$.

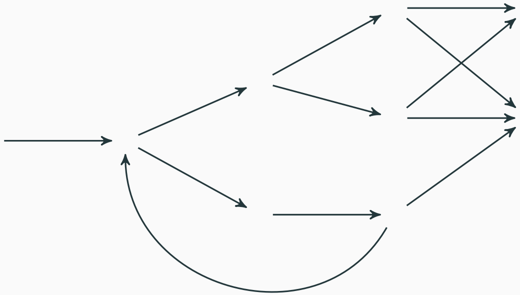
The flow of time



- Transitivity corresponds to **4**: $\mathsf{G}A \rightarrow \mathsf{G}\mathsf{G}A$.
- Asymmetry corresponds to nothing.
- Discreteness corresponds to $(\mathsf{F}(p \vee \neg p) \wedge A \wedge \mathsf{H}A) \rightarrow \mathsf{F}\mathsf{G}A$.
- Left-linearity corresponds to $\mathsf{F}\mathsf{P}A \rightarrow (\mathsf{F}A \vee A \vee \mathsf{P}A)$.
- Right-branchingness corresponds to nothing.

The flow of time

Relativistic time:



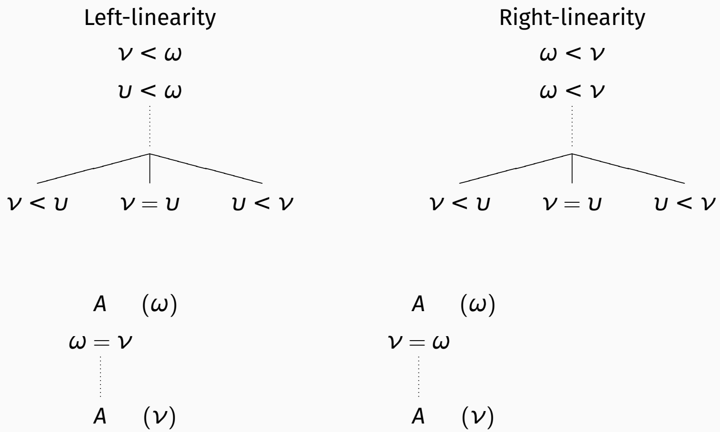
Trees

Trees

$\mathrm{G}A \quad (\omega)$	$\mathrm{F}A \quad (\omega)$	$\mathrm{H}A \quad (\omega)$	$\mathrm{P}A \quad (\omega)$
$\omega < \nu$	⋮	$\nu < \omega$	⋮
⋮	$\omega < \nu$	⋮	$\nu < \omega$
$A \quad (\nu)$	$A \quad (\nu)$	$A \quad (\nu)$	$A \quad (\nu)$
↑ old	↑ new	↑ old	↑ new

$\neg\,\mathrm{G}A \quad (\omega)$	$\neg\,\mathrm{F}A \quad (\omega)$	$\neg\,\mathrm{H}A \quad (\omega)$	$\neg\,\mathrm{P}A \quad (\omega)$
$\omega < \nu$	$\omega < \nu$	$\nu < \omega$	⋮
⋮	⋮	⋮	$\nu < \omega$
$\neg A \quad (\nu)$	$\neg A \quad (\nu)$	$\neg A \quad (\nu)$	$\neg A \quad (\nu)$
↑ new	↑ old	↑ new	↑ old

Trees


Left-linearity
ν < ω
υ < ω
ν < υ
ν = υ
υ < ν
Right-linearity
ω < ν
ω < ν
ν < υ
ν = υ
υ < ν
A (ω)
ω = ν
A (ν)
A (ω)
ν = ω
A (ν)